

The NII-Yamanashi-LRI Workshop

# Exhaustive analysis of the dynamics of Process Hitting through Answer Set Programming

Emna BEN ABDALLAH & Maxime FOSCHETTE

joint work Olivier ROUX & Morgan MAGNIN

MeForBio / IRCCyN / École Centrale de Nantes (Nantes, France)  
[emna.ben-abdallah@irc-cyn.ec-nantes.fr](mailto:emna.ben-abdallah@irc-cyn.ec-nantes.fr)

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## Context and Aims

**MeForBio** team:  
Algebraic modelling to study  
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### 3) What for?

searching of PH **properties** through ASP (Fixed points, reachability).

# Plan

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- ASP implementation
- Optimization of ASP implementation
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## 5) Conclusion & prospect

# Answer Set Programming

## ASP:

- Logic program written in language of AnsProlog\*
- Form of rules :

*head*  $\leftarrow$  *body*.  
 $L_0 \leftarrow L_1, \dots, L_m, \text{not } L_{m+1}, \dots, \text{not } L_n.$

with each  $L_i$  : literal in the sense of classical logic.

Rule's meaning:

If  $L_1, \dots, L_m$  are **true** and if  $L_{m+1}, \dots, L_n$  are **false**  
then  $L_0$  is **true**.

# Answer Set Programming

Special types of rules:

- **Constraint :**

$$\leftarrow L_1, \dots, L_m, \text{not } L_{m+1}, \dots, \text{not } L_n.$$

- **Fact :**

$$L_0.$$

- **Cardinality :**

$$\min\{L_0, \dots, L_j\} \max \leftarrow L_1, \dots, L_m, \text{not } L_{m+1}, \dots, \text{not } L_n.$$

# Answer Set Programming

## Example:

```
bird(X) ← lays_egg(X).  
mammal(X) ← engender(X).  
fly(X) ← bird(X), not mammal(X).  
lays_egg(tweety).
```

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## Solution:

```
bird(tweety) ← True.  
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```

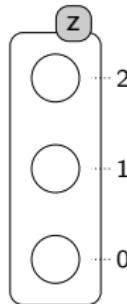
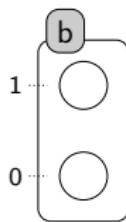
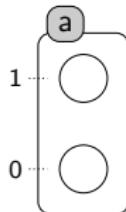
Answer: *fly(tweety)*, *bird(tweety)*.

# The Process Hitting modeling



**Sorts:** components  $a, b, z$

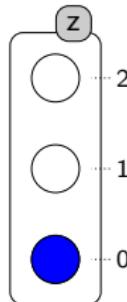
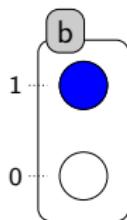
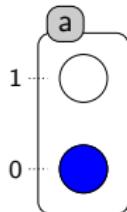
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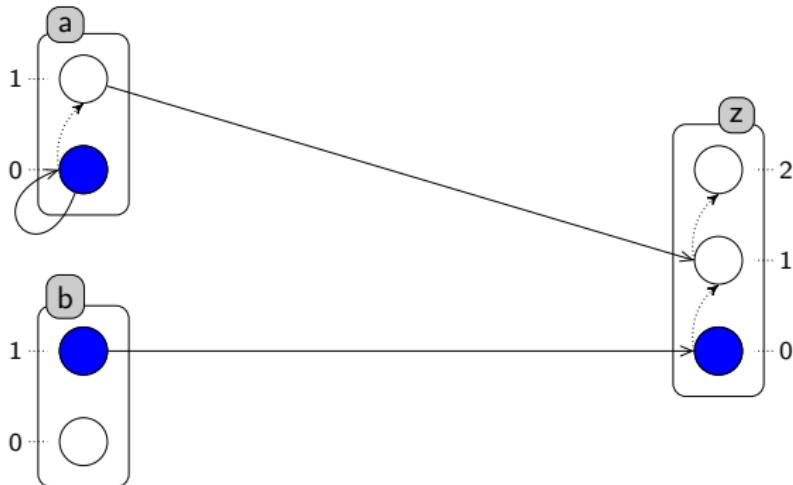


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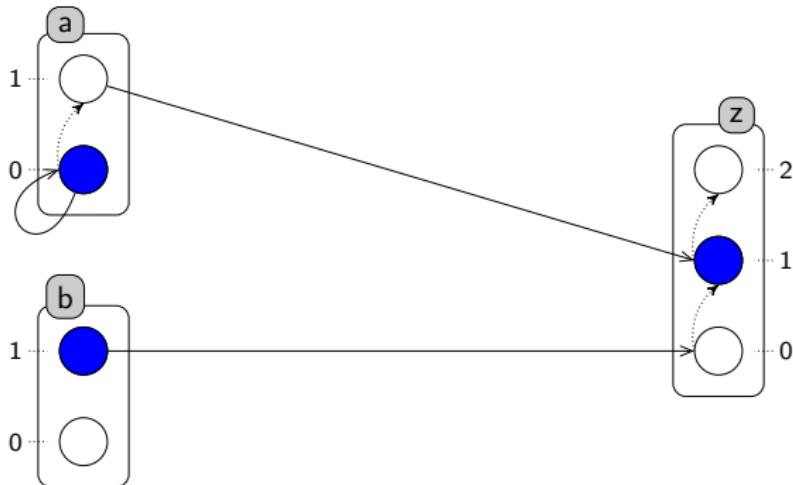
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**Actions:** dynamics  $b_1 \rightarrow z_0 \uparrow z_1, a_0 \rightarrow a_0 \uparrow a_1, a_1 \rightarrow z_1 \uparrow z_2$

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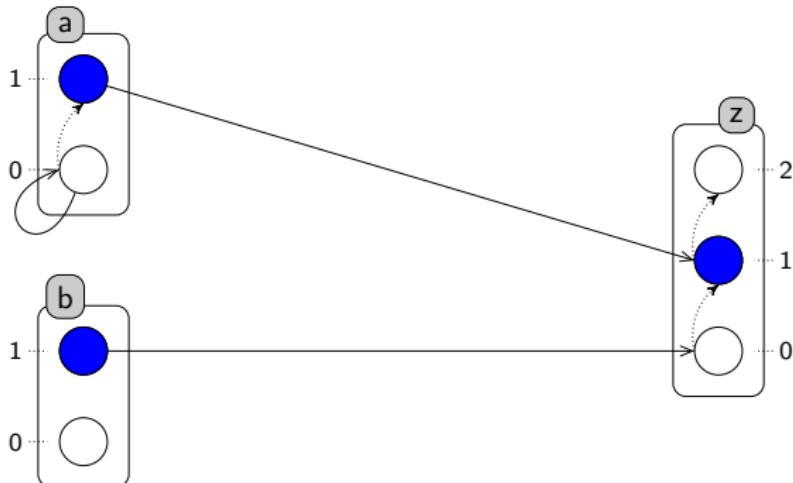
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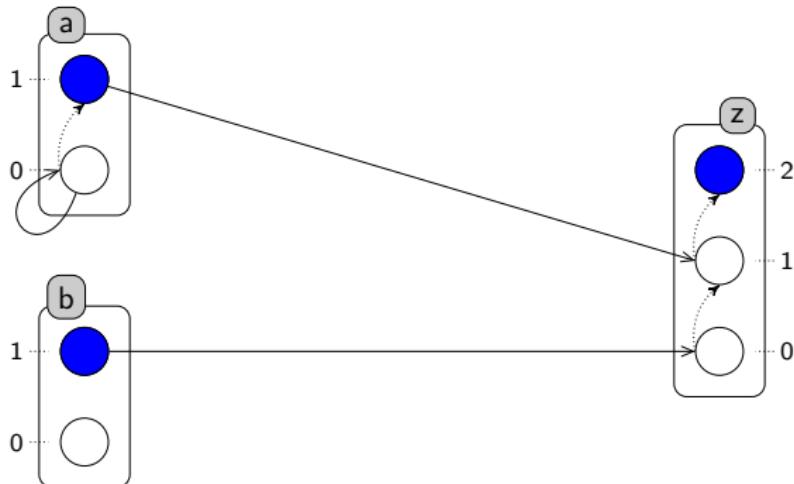
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# PH through ASP

## Network traduction:

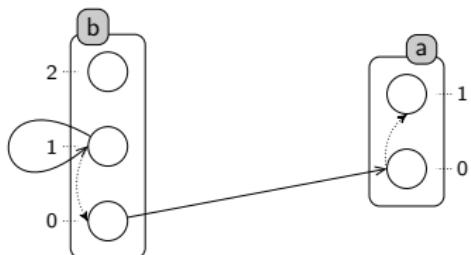
- **Sort**: `sort(A)` .
- **Process**: `process(A,I).`
- **Action**  $a_i \rightarrow b_j \uparrow b_k$  : `action(A,I,B,J,K)` .

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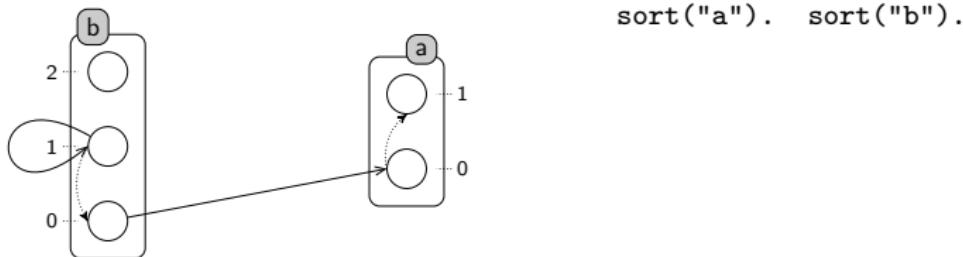


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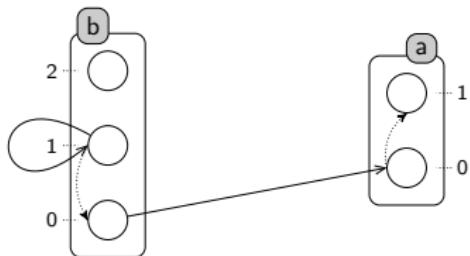


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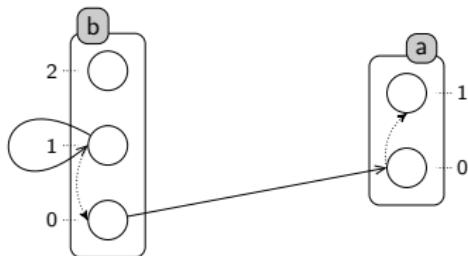
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sort("a"). sort("b").
process("a", 0..1).
process("b", 0..2).
```

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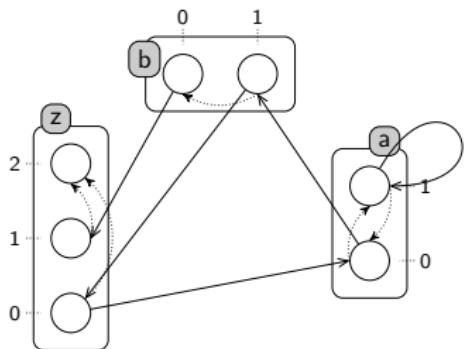
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sort("a") . sort("b") .
process("a", 0..1) .
process("b", 0..2) .
action("b", 0, "a", 0, 1) .
action("b", 1, "b", 1, 0) .

```

## Fixed Points

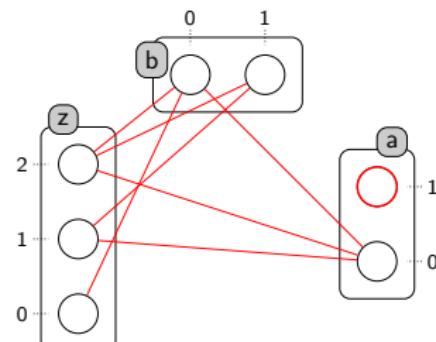
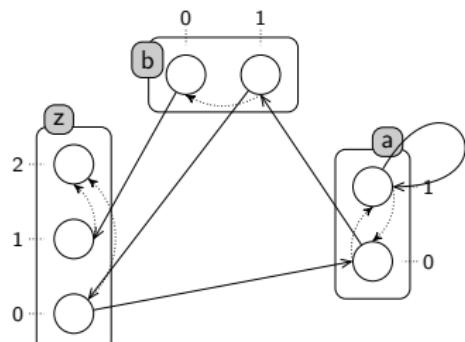
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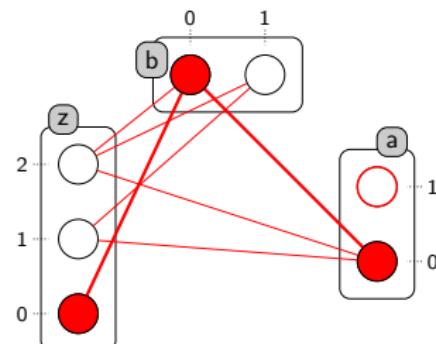
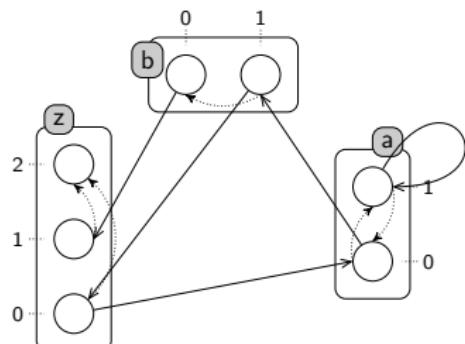
→ Hitless Graph



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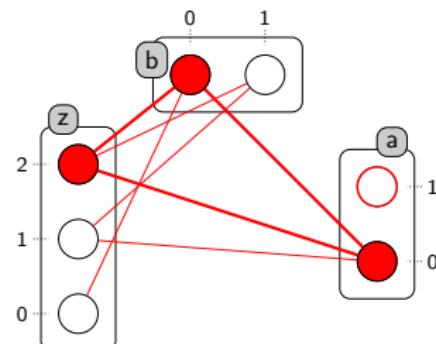
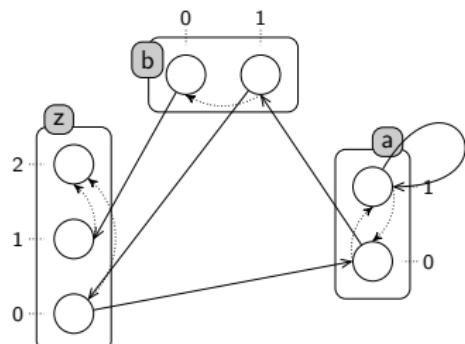
→ Hitless Graph → **n-clicks** = fixed points



# Fixed Points

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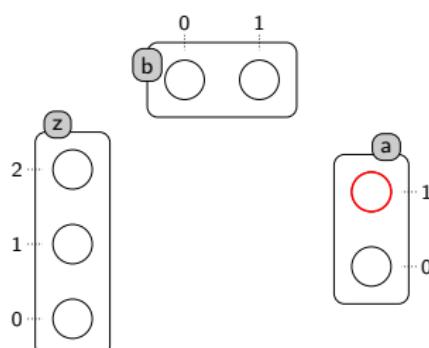
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# Implementation of the algorithm (N-Cliques)

Definition of hitless graph :

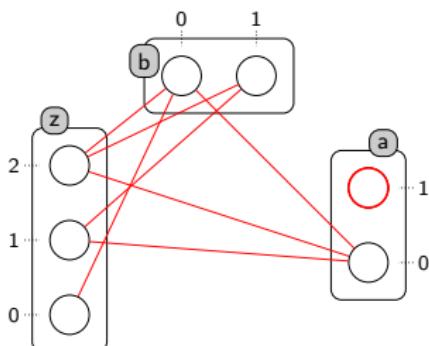
```
noAction(A,I,B,J) :- not hit(A,I,B,J), not hit(B,J,A,I), A!=B,  
                      shownProcess(A,I), shownProcess(B,J).
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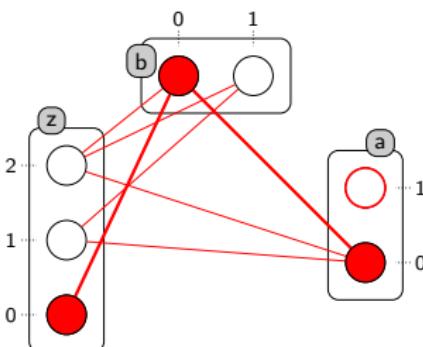
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Select processes :

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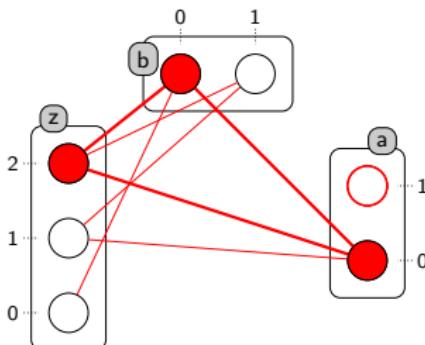
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Find Fixed points :

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noHit(A,I,B,J) :- noAction(A,I,B,J),  
selectProcess(A,I), selectProcess(B,J).  
noExistFixPoint :- 0 {noHit(A,I,B,J)} 0, selectProcess(A,I),  
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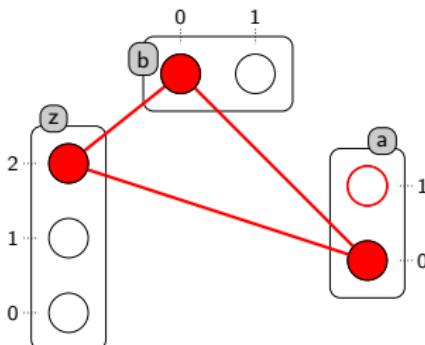
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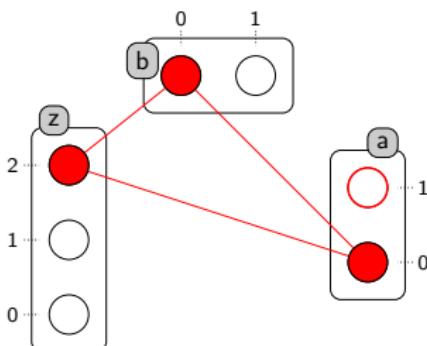


## Static analysis

Fixed point through ASP

### ASP program result:

Answer 1: fixProcess(a,0), fixProcess(b,0), fixProcess(z,2).



# Static analysis

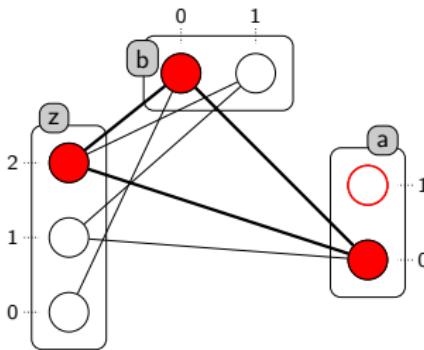
## Fixed point through ASP

### Optimization:

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noHit(A,I,B,J) :- noAction(A,I,B,J), selectProcess(A,I),
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-noExistFixPoint :- 0 {noHit(A,I,B,J)} 0, selectProcess(A,I),
                   selectProcess(B,J).
:- noExistFixPoint.
fixProcess(A,I) :- selectProcess(A,I).

```

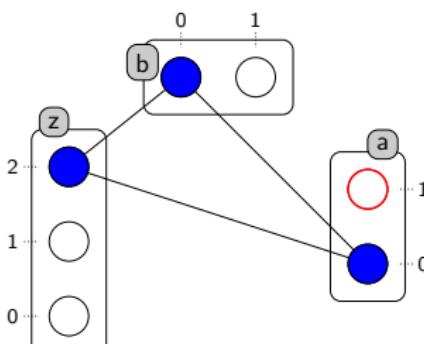


# Static analysis

Fixed point through ASP

## Optimization:

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:- hit(A,I,B,J), selectProcess(A,I), selectProcess(B,J), A!=B.
```



# Static analysis

## Fixed Point

### Comparison

Model	#sorts	#states	#fix-point	mthd1	mthd2	PINT
mvbrn	3	12	1	0.000s	0.000s	0.006s
ERBB	42	$2^{70}$	3	0.220s	0.000s	0.017s
tcrsig40	54	$2^{73}$	1	0.220s	0.020s	0.021s
tcrsig94	133	$2^{194}$	0	2.540s	0.060s	0.027s
egfr104	193	$2^{320}$	0	8.220s	0.140s	0.074s

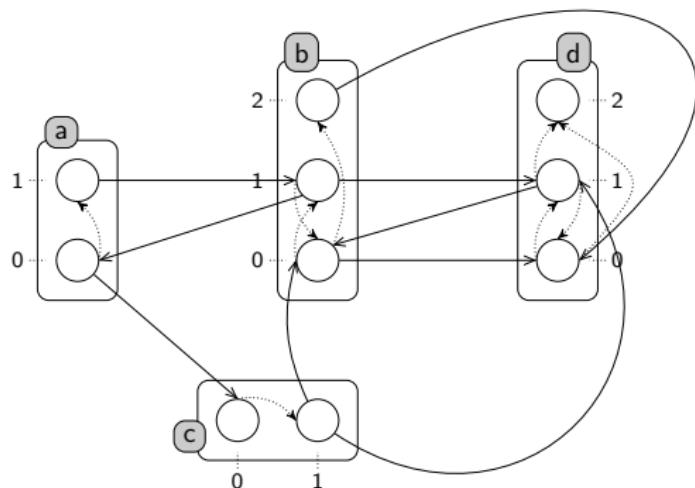
**Figure :** Execution time of ASP methods and PINT applied for biological networks with a desktop computer (core i5 and 4GB RAM).

**PINT** : a library developed to parse and study PH models.

# Dynamic analysis

## Reachability

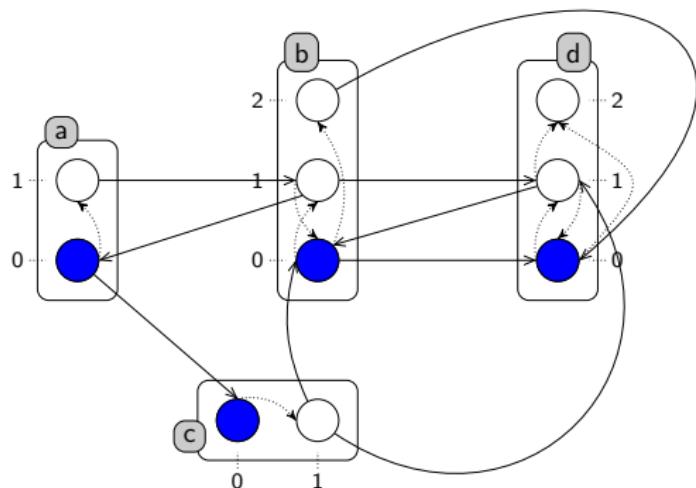
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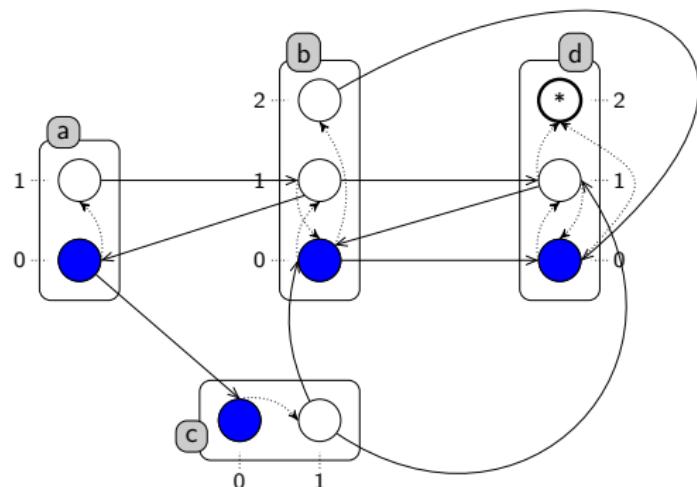
- Initial context

$$\langle a_0, b_0, c_0, z_0 \rangle$$

# Dynamic analysis

## Reachability

### Reachability of processes:



- Initial context

 $\langle a_0, b_0, c_0, z_0 \rangle$ 

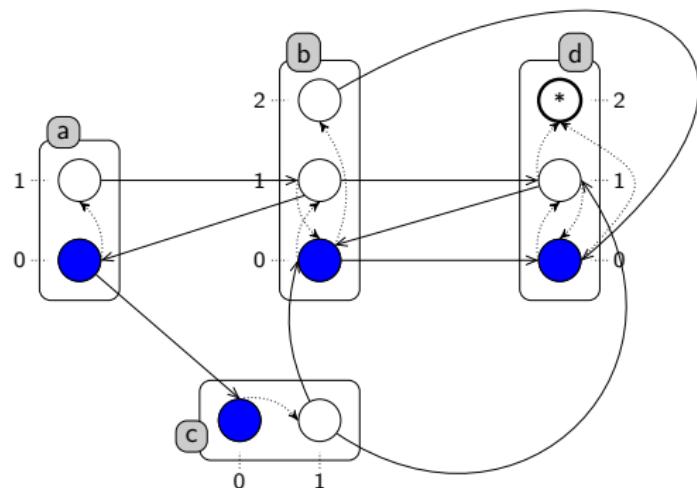
- Objectives

 $[ \triangleright d_2 ]$

# Dynamic analysis

## Reachability

### Reachability of processes:



- Initial context

$$\langle a_0, b_0, c_0, z_0 \rangle$$

- Objectives

$$[ \triangleright d_2 ]$$

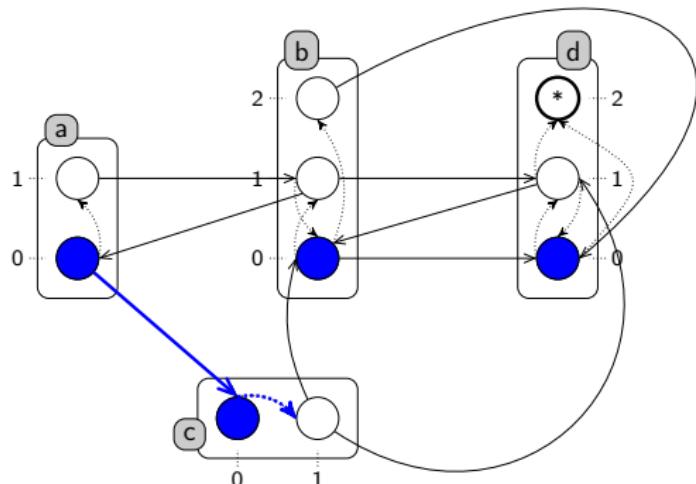
→ Concretization of the objective = scenario

$$a_0 \rightarrow c_0 \triangleright c_1 :: b_0 \rightarrow d_0 \triangleright d_1 :: c_1 \rightarrow b_0 \triangleright b_1 :: b_1 \rightarrow d_1 \triangleright d_2$$

# Dynamic analysis

## Reachability

### Reachability of processes:



- Initial context  
 $\langle a_0, b_0, c_0, z_0 \rangle$
- Objectives  
 $[ \triangleright d_2 ]$

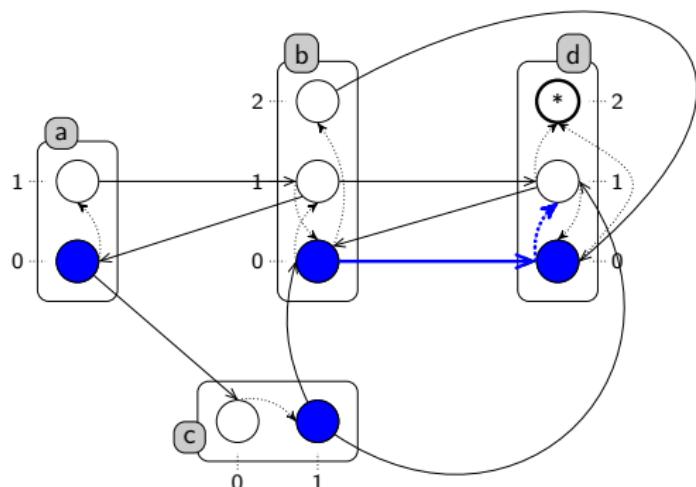
→ Concretization of the objective = scenario

$$\underline{a_0 \rightarrow c_0 \triangleright c_1 :: b_0 \rightarrow d_0 \triangleright d_1 :: c_1 \rightarrow b_0 \triangleright b_1 :: b_1 \rightarrow d_1 \triangleright d_2}$$

# Dynamic analysis

## Reachability

### Reachability of processes:



- Initial context

$$\langle a_0, b_0, c_0, z_0 \rangle$$

- Objectives

$$[ \triangleright d_2 ]$$

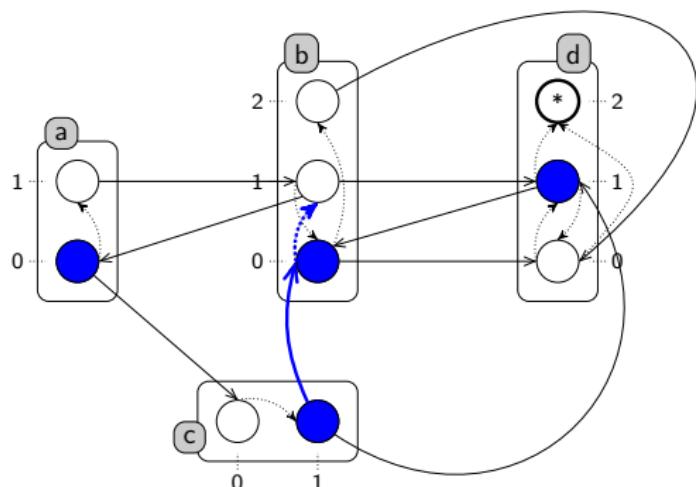
→ Concretization of the objective = scenario

$$a_0 \rightarrow c_0 \triangleright c_1 :: \underline{b_0 \rightarrow d_0 \triangleright d_1} :: c_1 \rightarrow b_0 \triangleright b_1 :: b_1 \rightarrow d_1 \triangleright d_2$$

# Dynamic analysis

## Reachability

### Reachability of processes:



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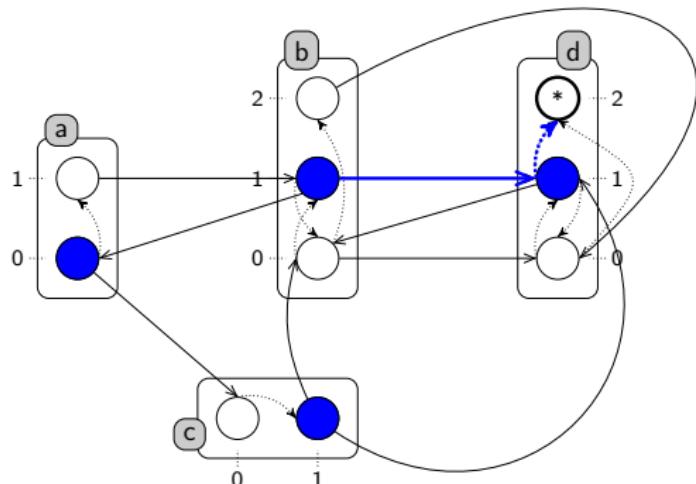
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$$a_0 \rightarrow c_0 \triangleright c_1 :: b_0 \rightarrow d_0 \triangleright d_1 :: \underline{c_1 \rightarrow b_0 \triangleright b_1} :: b_1 \rightarrow d_1 \triangleright d_2$$

# Dynamic analysis

## Reachability

### Reachability of processes:



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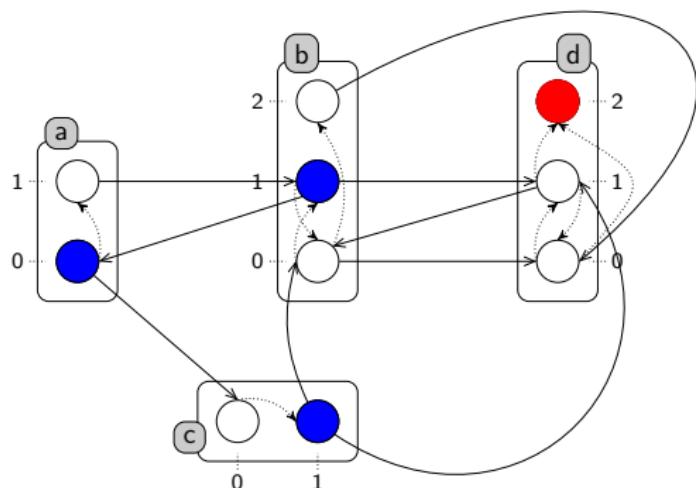
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# Dynamic analysis

## Reachability

### Reachability of processes:



- Initial context

$$\langle a_0, b_0, c_0, z_0 \rangle$$

- Objectives

$$[ \triangleright d_2 ]$$

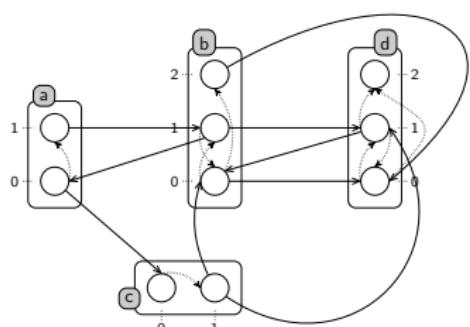
→ Concretization of the objective = scenario

$$a_0 \rightarrow c_0 \triangleright c_1 :: b_0 \rightarrow d_0 \triangleright d_1 :: c_1 \rightarrow b_0 \triangleright b_1 :: b_1 \rightarrow d_1 \triangleright d_2$$

# Dynamic analysis

Evolution through ASP

## Network evolution through ASP



# Dynamic analysis

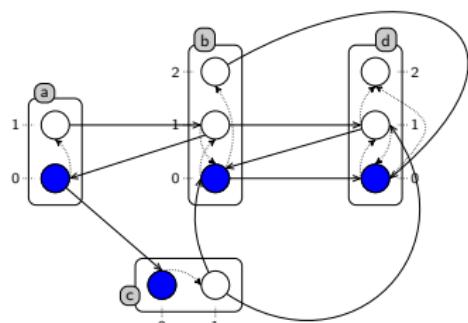
Evolution through ASP

## Network evolution through ASP

**Initializing :**

```
init(activeProcess("a",0)).
```

avec a: sorte, 0: indice du processus



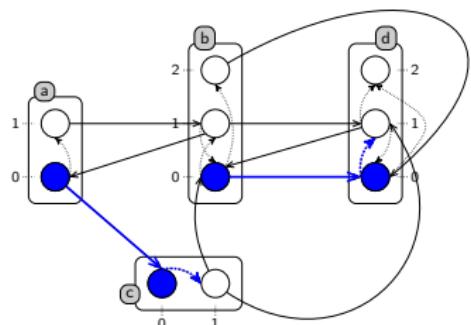
# Dynamic analysis

Evolution through ASP

## Network evolution through ASP

### Playable actions at step T :

```
playableAction(A,I,B,J,K,T) :- action(A,I,B,J,K),  
                                instate(activeProcess(A,I),T),  
                                instate(activeProcess(B,J),T), time(T).
```



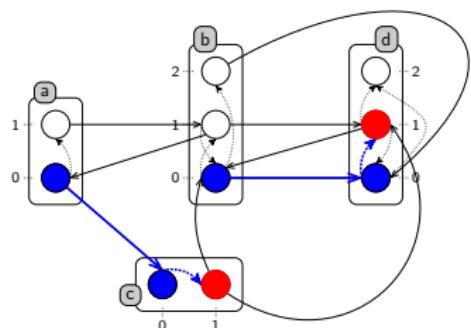
# Dynamic analysis

Evolution through ASP

## Network evolution through ASP

### Change active processes :

```
{activeFromTo(B,J,K,T)} :- playableAction(A,I,B,J,K,T),  
    J!=K, time(T).  
:- 2{ activeFromTo(B,J,K,T)}, time(T).
```



# Dynamic analysis

## Evolution through ASP

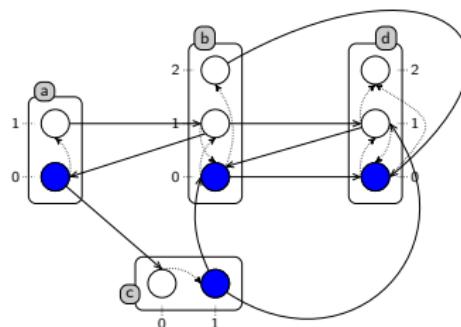
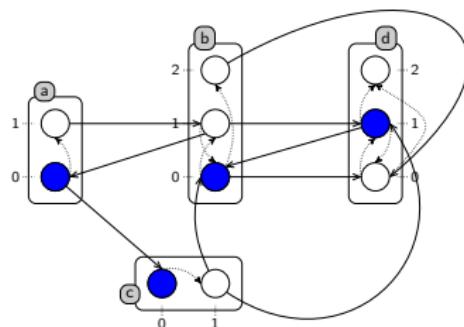
### Network evolution through ASP

#### Active processes at next step (T+1) :

```

instate(activeProcess(B,K),T+1) :- activeFromTo(B,J,K,T), time(T).
instate(activeProcess(A,I),T+1) :- instate(activeProcess(A,I),T),
                                activeFromTo(B,J,K,T), A!=B, time(T).

```



# Dynamic analysis

Evolution through ASP

## Network evolution through ASP

```
time(0..N).
```

### Results ( $N = 3$ ) :

Answer 1: activeFromTo("d",0,1,0) activeFromTo("c",0,1,1)  
activeFromTo("b",0,1,2).

Answer 2: activeFromTo("d",0,1,0) activeFromTo("b",0,2,1)

Answer 3: activeFromTo("c",0,1,0) activeFromTo("d",0,1,1)  
activeFromTo("d",1,0,2) activeFromTo("b",0,1,3)

...

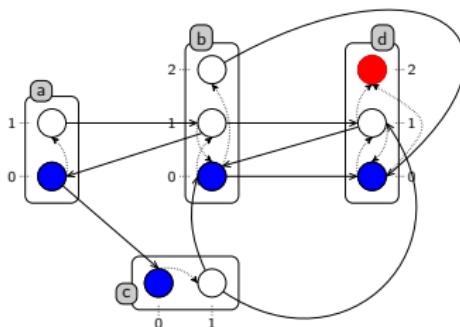
Answer 29: activeFromTo("c",0,1,0) activeFromTo("b",0,1,1)  
activeFromTo("a",0,1,2)

# Dynamic analysis

## Reachability through ASP

### Success reachability through ASP:

```
goal(activeProcess("d",2)).
```

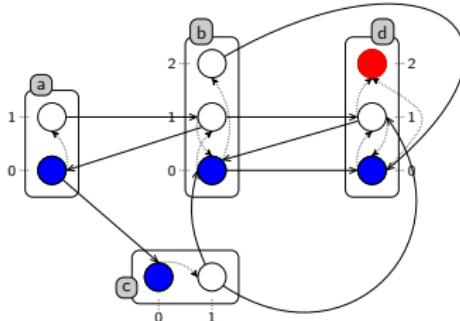


# Dynamic analysis

## Reachability through ASP

### Success reachability through ASP:

```
goal(activeProcess("d",2)).  
satisfiable(F,T) :- goal(F), instate(F,T).  
:- not satisfiableTot.
```

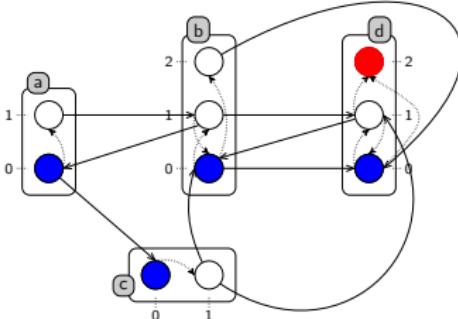


# Dynamic analysis

## Reachability through ASP

### Success reachability through ASP:

```
goal(activeProcess("d",2)).  
satisfiable(F,T) :- goal(F), instate(F,T).  
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time(0..N).
```

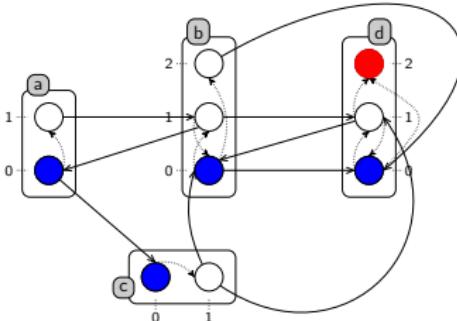


# Dynamic analysis

## Reachability through ASP

### Success reachability through ASP:

```
goal(activeProcess("d",2)).  
satisfiable(F,T) :- goal(F), instate(F,T).  
:- not satisfiableTot.  
time(0..N).  
predict N -> Inconvenient
```



# Dynamic analysis

## Reachability through ASP

**Results for ( $N = 2$ ) :**

UNSATISFIABLE

**Results for ( $N = 3$ ) :**

Answer 1: activeFromTo(c,0,1,0), activeFromTo(d,0,1,1),  
activeFromTo(b,0,1,2), activeFromTo(d,1,2,3).

Answer 2: activeFromTo("d",0,1,0) activeFromTo("c",0,1,1)  
activeFromTo("b",0,1,2) activeFromTo("d",1,2,3)

# Dynamic analysis

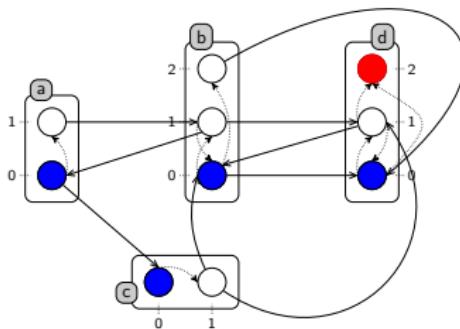
Reachability through ASP

## Success reachability through ASP iterative:

```

goal(activeProcess("d",2)).
#base
instate(F,0) :- init(F).
#cumulative t
playableAction(A, I, B, J, K,t), activeFromTo(B, J, K,t),
instate(activeProcess(A, I),t + 1)...
#volatile t
notSatisfiable(t) :- goal(F), not instate(F,t).
:- notSatisfiable(t).

```



# Dynamic analysis

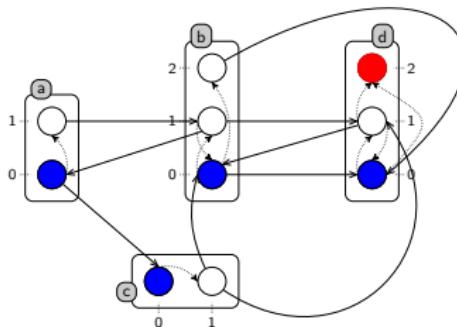
## Reachability through ASP

### Success reachability through ASP iterative:

#### Results:

Answer 1: activeFromTo(c,0,1,0), activeFromTo(d,0,1,1),  
activeFromTo(b,0,1,2), activeFromTo(d,1,2,3).

Answer 2: activeFromTo("d",0,1,0) activeFromTo("c",0,1,1)  
activeFromTo("b",0,1,2) activeFromTo("d",1,2,3)



# Dynamic analysis

## Reachability through ASP

### Comparison:

Initializing biological models components and the objectives.

Model	#sorts	#states	#steps	ASP	ASPi	PINT
Exemple	4	36	4	0.000s	0.000s	0.000s
ERBB	42	$2^{70}$	18	10.620s	5.020s	0.022s
tcrsig40	54	$2^{73}$	26	156.500s	127.250s	0.012s

Figure : Execution time of ASP methods ( CLINGO et ICLINGO ) and PINT applied for biological networks with a desktop computer (core i5 and 4GB RAM)

# Dynamic analysis

## Reachability through ASP

### **Comparison:**

Method of Rocca et al.:

- ASP
- CTL properties with model checking (AF, EF, AG...)
- Transitions graph

# Dynamic analysis

## Reachability through ASP

### Comparison:

Method of Rocca et al.:

- ASP
- CTL properties with model checking (AF, EF, AG...)
- Transitions graph

Comparaison of the property **EF**

prop = **EF** ( $I_0, goal$ )

## Dynamic analysis

Reachability through ASP

### Comparaison:

**Example:** Tail resorption of tadpole :

12 sorts, 42 process, 139 actions and 524.288 states.

$$\text{prop} = \text{EF}(l_0, goal)$$

# Dynamic analysis

## Reachability through ASP

### Comparaison:

**Example:** Tail resorption of tadpole :

12 sorts, 42 process, 139 actions and 524.288 states.

$$\text{prop} = \text{EF}(l_0, goal)$$

### Network traduction :

- Transition graph: 3min6s
- Process Hitting : 0.346s

# Dynamic analysis

## Reachability through ASP

### Comparaison:

**Example:** Tail resorption of tadpole :

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$$\text{prop} = \text{EF}(I_0, goal)$$

### Network traduction :

- Transition graph: 3min6s
- Process Hitting : 0.346s

### Property verification :

- Rocca et al. method : 7min17s
- our iterative method : 1.9s

# Conclusion & Prospects

- New dynamic analysis of Process Hitting models:
  - Fixed point
  - Network evolution
  - Reachability
- Prospects:
  - Adaptation on other models (PN, model of Thomas...)
  - Eliminating cycles
  - Search attractors
  - Reverse reachability ( $goal \rightarrow I_0?$ )

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**Thanks for your attention**