Réunion finale BioTempo

Introducing time parameters into Process Hitting with classes of priorities

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Joint work with: Loïc PAULEVÉ, Morgan MAGNIN, Olivier ROUX

Context and Aims

MeForBio team: Algebraic modeling to study large dynamical biological systems

→ Contribution: the Process Hitting framework [Paulevé *et al.* in Transactions on Computational Systems Biology, 2011] [Paulevé *et al.* in Mathematical Structures in Computer Science, 2012]

- · A restriction of synchronous automata networks
- Special form for the actions \Rightarrow more atomistic than Interaction Graphs
- Efficient reachability analysis
- \rightarrow Introduction of temporal features:

1) Stochastic parameters

[Paulevé et al. in Transactions on Computational Systems Biology, 2011]

2) Priorities

[Folschette et al. in Workshop on Interactions between Computer Science and Biology, 2013]

3) Neutralizing edges

The Process Hitting modeling

[Paulevé et al. in Transactions on Computational Systems Biology, 2011]



Sorts: components a, b, z

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Sorts: components *a*, *b*, *z* **Processes**: local states / levels of expression z_0 , z_1 , z_2

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Sorts: components *a*, *b*, *z* **Processes:** local states / levels of expression *z*₀, *z*₁, *z*₂ **States:** sets of active processes $\langle a_1, b_1, z_2 \rangle$ **Actions:** dynamics $b_1 \rightarrow z_0 \uparrow z_1$, $a_0 \rightarrow a_0 \uparrow a_1$, $a_1 \rightarrow z_1 \uparrow z_2$

Adding cooperations

[Paulevé et al. in Transactions on Computational Systems Biology, 2011]



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Cooperation between a_1 and b_1 : $a_1 \wedge b_1 \rightarrow z_0 \uparrow z_1$ Solution: a **cooperative sort** ab to express $a_1 \wedge b_1$

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Static analysis: successive reachability

[Paulevé et al. in Mathematical Structures in Computer Science, 2012]



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Successive reachability of processes:



 $\begin{array}{l} \rightarrow \text{ Concretization of the objective} = \text{scenario} \\ a_0 \rightarrow c_0 \mathrel{\sc r} c_1 :: b_0 \rightarrow d_0 \mathrel{\sc r} d_1 :: c_1 \rightarrow b_0 \mathrel{\sc r} b_1 :: b_1 \rightarrow d_1 \mathrel{\sc r} d_2 \end{array}$

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Over- and Under-approximations

[Paulevé et al. in Mathematical Structures in Computer Science, 2012]

Static analysis by abstractions:

- ightarrow Directly checking an objective sequence R is hard
- \rightarrow Rather check the approximations *P* and *Q*, where *P* \Rightarrow *R* \Rightarrow *Q*:



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Polynomial w.r.t. the number of sorts and exponential w.r.t. the number of processes in each sort

 \rightarrow Efficient for big models with few levels of expression



Sufficient condition:

- no cycle
- each objective has a solution





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- no cycle
- · each objective has a solution

R is true









Necessary condition:





Necessary condition:

There exists a traversal with no cycle

- objective \rightarrow follow one solution
- solution \rightarrow follow all processes
- process \rightarrow follow all objectives





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Inconclusive



The Process Hitting modeling

- Dynamic modeling with an atomistic point of view
 - \rightarrow Independent actions
 - \rightarrow Cooperation modeled with cooperative sorts
- Efficient static analysis
 - \rightarrow Reachability of a process can be computed in polynomial time in the number of sorts
 - \rightarrow Useful for the study of large biological models (up to hundreds of sorts)
- Results:

Model	Sorts	Procs	Actions	States	Biocham ¹	libddd ²	PINT
egfr20	35	196	670	264	[3s−∞]	[1s–150s]	0.007s
tcrsig40	54	156	301	273	[1s−∞]	[0.6s–∞]	0.004s
tcrsig94	133	448	1124	2 ¹⁹⁴	∞	∞	0.030s
egfr104	193	748	2356	2 ³²⁰	∞	∞	0.050s

¹ Inria Paris-Rocquencourt/Contraintes

² LIP6/Move

egfr20: [Epidermal Growth Factor Receptor, by Özgür Sahin et al.] egfr104: [Epidermal Growth Factor Receptor, by Regina Samaga et al.] tcrsig40: [T-Cell Receptor Signaling, by Steffen Klamt et al.] tcrsig94: [T-Cell Receptor Signaling, by Julio Saez-Rodriguez et al.]






























































Introducing time into PH with priorities o Stochastic Features

Stochastic Features

[Paulevé et al. in Transactions on Computational Systems Biology, 2011]

- Introduces time features
- Parameters: either (r, sa), or the firing interval [d; D].



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Introducing time into PH with priorities o Stochastic Features o Adding Stochasticity to the Metazoan Segmentation

Metazoan Segmentation with Stochastic Parameters

[Paulevé et al. in Transactions on Computational Systems Biology, 2011]



Introducing time into PH with priorities o Stochastic Features o Adding Stochasticity to the Metazoan Segmentation

Temporal Simulation [Paulevé (PhD thesis), 2011]



• Simulation with stochastic parameters:

- Other possible analysis: stochastic model checkers (PRISM)
 - ightarrow But combinatoric explosion: PRISM fails for more than 5 components

Pros and Cons of Stochastic Parameters

Pros:

- Introduction of temporal features
- Simulation in continuous time

Cons:

- Very hard to analyze
 - Either multiple runs (statistics)
 - Or model checkers (PRISM) but combinatorial explosion

Introduction of Classes of Priorities

- Each action is associated to a class of priority.
- An action cannot be played if another action of higher priority is playable.



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Metazoan Segmentation with Priorities



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Metazoan Segmentation with Priorities



Metazoan Segmentation in Canonical Form



Metazoan Segmentation in Canonical Form


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Introducing time into PH with priorities o Priorities o Adding Priorities to the Metazoan Segmentation

Metazoan Segmentation in Canonical Form



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Metazoan Segmentation in Canonical Form



 \rightarrow Same dynamics but only 2 priorities \rightarrow Priority 1 is only for cooperative sorts

Pros and Cons of Classes of Priorities

Pros:

- Better expressivity (same as Boolean Networks!)
- Efficient static analysis

Model	Sorts	Procs	Actions	States	libddd ¹	PINT
tcrsig94	133	448	1124	2 ¹⁹⁴	∞	0.008s - 0.060s

 1 LIP6/Move

tcrsig94: [T-Cell Receptor Signaling, by Julio Saez-Rodriguez *et al.*] (Here with prioritized cooperative sorts)

Cons:

No accumulation phenomenons

Contrary to stochastic simulation

- Translation to canonical form is exponential
 - ightarrow For each action of priority *n*, exponential in the number of actions of priority $[\![2;n-1]\!]$

Neutralizing Edges



- - $a_0
 ightarrow b_0
 ightarrow b_1$ is playable
- ightarrow Here, only one possible behavior

Neutralizing Edges



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Metazoan Segmentation with Neutralizing Edges



Pros and Cons of Neutralizing Edges

Pros:

- Same expressivity than Priorities
 - $\rightarrow~$ Can be translated to the canonical form
- Finer preemption relations
 - \rightarrow Easier modeling in some cases
- Sparser constraints
 - $\rightarrow~$ More efficient translation to canonical form

Summary & Conclusion

Process Hitting: an atomistic modeling with powerful static analysis

- 1. Stochastic parameters:
 - · To model systems with chronometric features
 - Continuous time
 - But hard to analyze
- 2. Classes of priorities:
 - Allows to reproduce the same behaviors
 - Efficient static analysis
 - But the translation to canonical form faces combinatorial explosion
- 3. Neutralizing edges:
 - Alternative to priorities
 - · Closer to reality in some cases
 - Lighter translation to canonical form

Thank you

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Introducing time into PH with priorities o Adding Priorities

Adding cooperations

[Paulevé et al. in Transactions on Computational Systems Biology, 2011]



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Drawback: Cooperations are too "loose" to be as expressive as ADN. $\langle a_0, b_0, ab_{00}, z_0 \rangle$

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$$\langle \mathbf{a}_0, \mathbf{b}_0, \mathbf{a}\mathbf{b}_{00}, \mathbf{z}_0 \rangle \rightarrow \langle \mathbf{a}_1, \mathbf{b}_0, \mathbf{a}\mathbf{b}_{00}, \mathbf{z}_0 \rangle \rightarrow \langle \mathbf{a}_1, \mathbf{b}_0, \mathbf{a}\mathbf{b}_{10}, \mathbf{z}_0 \rangle$$

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$$\langle \mathsf{a}_0, \mathsf{b}_0, \mathsf{a}\mathsf{b}_{00}, \mathsf{z}_0 \rangle \rightarrow \langle \mathsf{a}_1, \mathsf{b}_0, \mathsf{a}\mathsf{b}_{00}, \mathsf{z}_0 \rangle \rightarrow \langle \mathsf{a}_1, \mathsf{b}_0, \mathsf{a}\mathsf{b}_{10}, \mathsf{z}_0 \rangle \rightarrow \langle \mathsf{a}_0, \mathsf{b}_0, \mathsf{a}\mathsf{b}_{10}, \mathsf{z}_0 \rangle$$

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[Folschette et al. in_Workshop on Interactions between Computer Science and Biology, 2013]



Drawback: Cooperations are too "loose" to be as expressive as ADN.

$$\begin{array}{l} \langle \mathsf{a}_0, \mathsf{b}_0, \mathsf{a}\mathsf{b}_{00}, \mathsf{z}_0 \rangle \to \langle \mathsf{a}_1, \mathsf{b}_0, \mathsf{a}\mathsf{b}_{00}, \mathsf{z}_0 \rangle \to \langle \mathsf{a}_1, \mathsf{b}_0, \mathsf{a}\mathsf{b}_{10}, \mathsf{z}_0 \rangle \to \langle \mathsf{a}_0, \mathsf{b}_1, \mathsf{a}\mathsf{b}_{10}, \mathsf{z}_0 \rangle \\ \to \langle \mathsf{a}_0, \mathsf{b}_1, \mathsf{a}\mathsf{b}_{10}, \mathsf{z}_0 \rangle \to \langle \mathsf{a}_0, \mathsf{b}_1, \mathsf{a}\mathsf{b}_{11}, \mathsf{z}_1 \rangle \to \langle \mathsf{a}_0, \mathsf{b}_1, \mathsf{a}\mathsf{b}_{11}, \mathsf{z}_1 \rangle \end{array}$$

The cooperativity should be: $a_1 \wedge b_1$ simultaneously *i.e.* "in the same state" but the model behaves like: $P(a_1) \wedge P(b_1)$ with P = "previously"

[Folschette et al. in Workshop on Interactions between Computer Science and Biology, 2013]



- Prioritise actions updating cooperative sorts (non-biological actions)
- All other actions remain unprioritised (evolutions with delays)

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[Folschette et al. in Workshop on Interactions between Computer Science and Biology, 2013]



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